Naip Moro

July 2017

# 1 Introduction

In chapter 7 of *Laws of Form* Spencer-Brown extends the scope of his basic equations to expressions with any finite number of variables. Some of his arguments, when he provides them, are rigorous; others are mere sketches, and some possible generalizations are left unmentioned. This paper will present fully rigorous proofs of the propositions.

Below is a list of axioms and theorems referenced in subsequent proofs:

$$\overline{pr} |\overline{qr}| = \overline{p} |\overline{q}| r \tag{J2}$$

$$pr | \overline{qr} = \overline{p} | \overline{q} | r | \tag{J2.1}$$

$$a = a \tag{C1}$$

$$\begin{array}{c} ab | b = a | b \\ \hline \end{array} \tag{C2}$$

$$a|b|c| = ac|b|c| \tag{C7}$$

$$\overrightarrow{a} \overrightarrow{br} \overrightarrow{cr} = \overrightarrow{a} \overrightarrow{b} \overrightarrow{c} \overrightarrow{a} \overrightarrow{r}$$
(C8)

$$\overrightarrow{a|r|}\overrightarrow{b|r|}\overrightarrow{x|r}\overrightarrow{y|r|} = \overrightarrow{r|ab}\overrightarrow{rxy}$$
(C9)

$$\overline{\overline{a|r|}} \overline{\overline{x|r|}} = \overline{\overline{r|a|rx|}}$$
(C9.1)

## 2 General theorems

Spencer-Brown begins the chapter by sketching an inductive generalization of J2. Here is the proof in full.

Theorem (J2\*).

$$\overline{\overline{a_1}} \overline{a_2} \dots \overline{a_n} r = \overline{\overline{a_1 r}} \overline{a_2 r} \dots \overline{a_n r}$$

*Proof.* The proof proceeds by induction on n. The base case is J2, where n = 2. Let the induction hypothesis (J2h) be:

$$\overline{a_1}$$
  $\overline{a_2}$  ...  $\overline{a_n}$   $r = \overline{a_1 r}$   $\overline{a_2 r}$  ...  $\overline{a_n r}$ 

The induction step:

$$\overline{a_1} \overline{a_2} \dots \overline{a_n} \overline{a_{n+1}} r$$

$$= \overline{\overline{a_1} \overline{a_2} \dots \overline{a_n}} \overline{a_{n+1}} r$$
(C1)

$$= \overline{a_1} \overline{a_2} \dots \overline{a_n} |r| \overline{a_{n+1}r}|$$
(J2)

$$= \overline{\overline{a_1 r} \overline{a_2 r} \dots \overline{a_n r}} \left\| \overline{a_{n+1} r} \right\|$$
(J2h)

$$= \overline{a_1 r} \overline{a_2 r} \dots \overline{a_n r} \overline{a_{n+1} r} |$$
(C1)

*Alternate proof.* A very similar and equally short proof, using the same induction hypothesis as above. The induction step:

$$\begin{array}{c}
\overline{a_{1}} \\
\overline{a_{2}} \\
\ldots \\
\overline{a_{n}} \\
\overline{a_{2}} \\
\ldots \\
\overline{a_{n}} \\
\overline{a_{n+1}} \\
\end{array} \\
r$$
(C1)

$$= \underline{\overline{a_1r} \ \overline{a_2r} \ \dots \ \overline{a_n} \ \overline{a_{n+1}} \ r}$$
(J2h)

$$= \overline{a_1 r} \overline{a_2 r} \dots \overline{a_n r} \overline{a_{n+1} r} |||$$
(J2)

$$= \overline{a_1 r} \overline{a_2 r} \dots \overline{a_n r} \overline{a_{n+1} r} |$$
(C1)

Before continuing, I prove a useful generalization of corollary J2.1.

### Theorem (J2.1\*).

$$\overline{a_1r} \ \overline{a_2r} \ \dots \ \overline{a_nr} = \overline{\overline{a_1} \ \overline{a_2} \ \dots \ \overline{a_n}} r$$

Proof.

$$\overline{a_1 r} \overline{a_2 r} \dots \overline{a_n r}$$

$$= \overline{\overline{a_1 r} \overline{a_2 r} \dots \overline{a_n r}}$$
(C1)

$$= \overline{a_1} a_2 \dots a_n |r| \tag{J2*}$$

Spencer-Brown states the generalizations of C8 and C9 but omits the proofs, merely noting that they are similar to  $J2^*$ .

#### Theorem (C8\*).

$$\overline{a} \overline{b_1 r} \overline{b_2 r} \dots \overline{b_n r} = \overline{a} \overline{b_1} \overline{b_2} \dots \overline{b_n} |\overline{a} \overline{r}|$$

*Proof.* The proof proceeds by induction on n. The base case is C8, where n = 2. Let the induction hypothesis (C8h) be:

$$\overline{a} \overline{b_1 r} \overline{b_2 r} \dots \overline{b_n r} = \overline{a} \overline{b_1} \overline{b_2} \dots \overline{b_n} \overline{a} \overline{r}$$

The induction step:

$$\overline{a} \overline{b_1 r} \overline{b_2 r} \dots \overline{b_n r} \overline{b_{n+1} r}$$

$$= \overline{\overline{a}} \overline{\overline{b_1 r}} \overline{\overline{b_2 r}} \dots \overline{\overline{b_n r}} \overline{\overline{b_{n+1} r}}$$

$$= \overline{\overline{a}} \overline{\overline{b_1 r}} \overline{\overline{b_2 r}} \dots \overline{\overline{b_n b_{n+1} r}}$$
(C1)
(C3)
(C3)
(C3)

$$= \frac{a|b_1r|||b_2|\dots b_n|b_{n+1}||a|b_1r|||r||}{\overline{b_2}|\dots \overline{b_n}|\overline{b_{n+1}}||\overline{r}||}$$
(C8h)  
$$= \frac{\overline{b_2}|\dots \overline{b_n}|\overline{b_{n+1}}||\overline{r}||\overline{a}|\overline{b_1r}|||}{\overline{a}|\overline{b_1r}|||}$$
(J2.1)

$$= \overline{\overline{b_2} \dots \overline{b_n} \overline{b_{n+1}}} r \overline{a} \overline{b_1 r}$$
 (C1 twice)

$$= \overline{\overline{b_2} \dots \overline{b_n} \overline{b_{n+1}}} \overline{b_1} r \overline{a}$$
(J2.1)

$$= \overline{\overline{b_1}} \overline{b_2} \dots \overline{b_n} \overline{b_{n+1}} r |\overline{a}|$$
(C1)

$$= \overline{b_1} \overline{b_2} \dots \overline{b_n} \overline{b_{n+1}} \overline{r} \| \overline{a} \|$$
(C1)

$$= \overline{a|b_1|b_2|\dots b_n|\overline{b_{n+1}}|} \overline{a|r|}$$
(J2)

$$= \overline{a} |b_1| |b_2| \dots |b_n| |b_{n+1}| |\overline{a}| \overline{r}|$$
(C1)

 $\mathrm{J2.1}^*$  allows for a quicker direct proof.

=

=

Alternate proof.

$$\overline{a} \overline{b_1 r} \overline{b_2 r} \dots \overline{b_n r}$$

$$= \overline{a} \overline{\overline{b_1} \overline{b_2} \dots \overline{b_n} r}$$
(J2.1\*)

$$= \overline{a} \overline{b_1} \overline{b_2} \dots \overline{b_n} \overline{r}$$
(C1)

$$= \overline{\overline{a} | \overline{b_1} | \overline{b_2} | \dots \overline{b_n} |} \overline{\overline{a} | \overline{r} |}$$
(J2)

$$= \overline{a} \overline{b_1} \overline{b_2} \dots \overline{b_n} || \overline{a} \overline{r} ||$$
(C1)

#### Theorem (C9\*).

$$\overline{\overline{a_1 | r|}} \overline{\overline{a_2 | r|}} \dots \overline{\overline{a_n | r|}} \overline{\overline{x_1 | r|}} \overline{\overline{x_2 | r|}} \dots \overline{\overline{x_m | r|}}$$
$$= \overline{\overline{r}} \overline{a_1 a_2 \dots a_n} \overline{\overline{rx_1 x_2 \dots x_m}}$$

Proof.

$$\overline{a_{1} | r|} \overline{a_{2} | r|} \dots \overline{a_{n} | r|} \overline{\overline{x_{1} | r|}} \overline{\overline{x_{2} | r|}} \dots \overline{\overline{x_{m} | r|}}$$

$$= \overline{\overline{a_{1} | \overline{a_{2} | \dots \overline{a_{n}} | r|}} \overline{\overline{x_{1} | \overline{x_{2} | \dots \overline{x_{m}} | r|}}$$

$$(J2.1^{*} twice)$$

$$= \overline{\overline{a_{1} a_{2} \dots a_{n} | r|} \overline{\overline{x_{1} x_{2} \dots x_{m} | r|}}$$

$$(C1 n+m times)$$

$$= \overline{r | a_{1} a_{2} \dots a_{n} | \overline{rx_{1} x_{2} \dots x_{m} | r|}}$$

$$(C9.1)$$

Next we prove a generalizion of C2.

#### Theorem $(C2^*)$ .

$$\boxed{\boxed{a_n b} \dots a_2} a_1 b = \boxed{\boxed{a_n} \dots a_2} a_1 b$$

*Proof.* The proof proceeds by induction on n. The base case is C2, where n = 1. Let the induction hypothesis be:

$$\boxed{\boxed{a_n b} \dots a_2} a_1 b = \boxed{\boxed{a_n} \dots a_2} a_1 b$$

Substitute  $\overline{a_{n+1}b} \mid a_n$  for  $a_n$  . The induction step then follows immediately:

$$\boxed{\boxed{a_{n+1}b} \ a_n b} \dots \ a_2 \ a_1 \ b = \boxed{\boxed{a_{n+1}b} \ a_n} \dots \ a_2 \ a_1 \ b$$

Spencer-Brown does not mention a generalized C7. Here is one possible version.

**Theorem** (C7\*). Let n be a positive even number. Then for all such n the following pair of equations holds:

(i) 
$$\overline{a_n} \dots a_2 a_1 = \overline{a_n} a_{n-1} \dots a_3 a_1 \dots \overline{a_4} a_3 a_1 \overline{a_2} a_1$$
  
(ii)  $\overline{a_{n+1}} a_n \dots a_2 a_1 = \overline{a_{n+1}} a_{n-1} \dots a_3 a_1 \overline{a_n} a_{n-1} \dots a_3 a_1 \dots \overline{a_4} a_3 a_1 \overline{a_2} a_1$ 

*Proof.* Let equation (i) be the induction hypothesis. The base case is the identity  $\overline{a_2} |a_1| = \overline{a_2} |a_1|$ , where n = 2. Now substitute  $\overline{a_{n+1}} |a_n|$  for  $a_n$ . Then,

$$\begin{array}{c}
\hline \hline a_{n+1} \ a_n \ \dots \ a_2 \ a_1 \\
= \ \hline a_{n+1} \ a_n \ a_{n-1} \dots \ a_3 a_1 \\
= \ \hline a_{n+1} \ a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_3 a_1 \\
\hline a_n \ a_{n-1} \dots \ a_{n-1} \dots \ a_{n-1} \\
\hline a_n \ a_{n-1} \dots \ a_{n-1} \dots \ a_{n-1} \\
\hline a_n \ a_{n-1} \dots \ a_{n-1} \dots \ a_{n-1} \\
\hline a_n \ a_{n-1} \dots \ a_{n-1} \dots \ a_{n-1} \\
\hline a_n \ a_{n-1} \dots \ a_{n-1} \dots \ a_{n-1} \\
\hline a_n \ a_{n-1} \dots \ a_{n-1} \dots \ a_{n-1} \\
\hline a_n \ a_{n-1} \dots \ a_{n-1} \dots \ a_{n-1} \\
\hline a_n \ a_{n-1} \dots \ a_{n-1} \dots \ a_{n-1} \\
\hline a_n \ a_{n-1} \dots \ a_{n-1} \dots \ a_{n-1} \dots \ a_{n-1} \\
\hline a_n \ a_{n-1} \dots \ a_{n-1} \dots \ a_{n-1} \dots \ a_{n-1} \\
\hline a_n \ a_{n-1} \dots \ a_{n-1} \dots$$

proving the implication from (i) to (ii). In equation (ii) substitute  $a_{n+2} |a_{n+1}|$  for  $a_{n+1}$ . Then,

$$\overline{\overline{a_{n+2}} a_{n+1}} \dots a_2 a_1 = \overline{a_{n+2}} a_{n+1} \dots a_3 a_1 \dots \overline{a_4} a_3 a_1 \overline{a_2} a_1$$
(ii)

proving (i) for the succeeding even number. This proves the proposition for all  $n \ge 2$ , and hence for all echelons of depth greater than or equal to 2.

**Theorem (T14).** Any expression can be reduced to an equivalent expression not more than two crosses deep. Specifically, any expression E is equivalent to  $\overline{a_1} \ b_1 \ \overline{a_2} \ b_2 \ \dots \ \overline{a_n} \ b_n \ \overline{c_1} \ \overline{c_2} \ \dots \ \overline{c_m} \ d$  where  $a_i, b_i, c_i, d$  are composed (at most) of juxtapositions of variables and the two constants,  $\neg$  and  $\neg$ .

*Proof.* Repeated applications of  $C7^*$  to any expression demonstrates the theorem. Spencer-Brown uses C7 (not having proven a generalization), but it comes to the same thing.

The final theorem follows Spencer-Brown closely.

**Theorem** (T15). Given any expression E and any variable v, E can be reduced to an equivalent expression containing not more than two appearances of v.

*Proof.* In the case where v is not in E, the theorem is trivially true, since  $E = \overline{v|v|} E$  by **J1**. So let us suppose that v appears in E. Using **C7\*** as many times as necessary, we rewrite E:

$$E = \overline{va_1} \ b_1 \ \overline{va_2} \ b_2 \ \dots \ \overline{va_n} \ b_n \ \overline{vc_1} \ \overline{vc_2} \ \dots \ \overline{vc_m} \ d$$

where  $a_i, b_i, c_i$ , and d are expressions free of v. Then, by n applications of **C8.1**,

$$E = \overrightarrow{v} \begin{vmatrix} b_1 & \overrightarrow{a_1} & b_1 & \overrightarrow{v} & b_2 & \overrightarrow{a_2} & b_2 & \cdots & \overrightarrow{v} & b_n & \overrightarrow{a_n} & b_n & \overrightarrow{vc_1} & \overrightarrow{vc_2} & \cdots & \overrightarrow{vc_m} & d \\ = \overrightarrow{v} \begin{vmatrix} b_1 & \overrightarrow{v} & b_2 & \cdots & \overrightarrow{v} & b_n & \overrightarrow{vc_1} & \overrightarrow{vc_2} & \cdots & \overrightarrow{vc_m} & f \\ (\text{where } f = \overrightarrow{a_1} & b_1 & \overrightarrow{a_2} & b_2 & \cdots & \overrightarrow{a_n} & b_n & d \text{ is free of } v.) \\ = \overline{\overrightarrow{b_1} & \overrightarrow{b_2} & \cdots & \overrightarrow{b_n}} \begin{vmatrix} \overrightarrow{v} & \overrightarrow{c_1} & \overrightarrow{c_2} & \cdots & \overrightarrow{c_m} & v \\ \hline{c_1} & \overrightarrow{c_2} & \cdots & \overrightarrow{c_m}} \begin{vmatrix} v & f & (J2.1^* \text{ twice}) \end{vmatrix}$$